

## 12.2 Representing Boolean Functions

### Literals

A *literal* is a Boolean variable or its complement.

### Minterm

A *minterm* of the Boolean variables  $x_1, x_2, \dots, x_n$  is a Boolean product  $y_1 y_2 \dots y_n$ , where  $y_i = x_i$  or  $y_i = \bar{x}_i$ . Hence, a minterm is a product of  $n$  literals, with one literal for each variable.

### Disjunctive Normal Form

The *disjunctive normal form* (DNF) of a *degree- $n$*  Boolean function  $f$  is the unique sum of minterms of the variables  $x_1, \dots, x_n$  that represents  $f$ .

### Maxterm

A *maxterm* of the Boolean variables  $x_1, x_2, \dots, x_n$  is a Boolean sum  $y_1 + y_2 + \dots + y_n$ , where  $y_i = x_i$  or  $y_i = \bar{x}_i$ . Hence, a maxterm is a sum of  $n$  literals, with one literal for each variable.

### Conjunctive Normal Form

The *conjunctive normal form* (CNF) of a *degree- $n$*  Boolean function  $f$  is the unique product of maxterms of the variables  $x_1, \dots, x_n$  that represents  $f$ .

### Functional Completeness

Since every Boolean function can be expressed in terms of  $\cdot, +, \bar{\phantom{x}}$ , we say that the set of operators  $\{\cdot, +, \bar{\phantom{x}}\}$  is *functionally complete*.  $\{+, \bar{\phantom{x}}\}$  and  $\{\cdot, \bar{\phantom{x}}\}$  are also functionally complete.  $\{\cdot, \bar{\phantom{x}}\}$  can be represented by the NAND operator  $|$  and  $\{+, \bar{\phantom{x}}\}$  by the NOR operator  $\downarrow$ . Therefore  $\{| \}$  and  $\{\downarrow\}$  are also functionally complete.

### 12.2 pg. 822 # 1

Find a Boolean product of the Boolean variables  $x, y,$ , and  $z$ , or their complements, that has the value 1 if and only if

a)  $x = y = 0, z = 1$

Since  $x = y = 0$ , then  $\bar{x} = \bar{y} = 1$ . The boolean product is simply  $\bar{x}\bar{y}z$ .

b)  $x = 0, y = 1, z = 0$

$$\bar{x}y\bar{z}$$

c)  $x = 0, y = z = 1$

$$\bar{x}yz$$

d)  $x = y = z = 0$

$\bar{x}\bar{y}\bar{z}$

**12.2 pg. 822 # 3**

Find the sum-of-products expansions of these Boolean functions.

a)  $F(x, y, z) = x + y + z$

Create the table for the function and take the Boolean sum of the minterm where the function evaluates to 1.

$x$	$y$	$z$	$x + y + z$
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

The DNF is  $\bar{x}\bar{y}z + \bar{x}y\bar{z} + \bar{x}yz + x\bar{y}\bar{z} + x\bar{y}z + xy\bar{z} + xyz$ 

b)  $F(x, y, z) = (x + z)y$

$x$	$y$	$z$	$x + z$	$(x + z)y$
0	0	0	0	0
0	0	1	1	0
0	1	0	0	0
0	1	1	1	1
1	0	0	1	0
1	0	1	1	0
1	1	0	1	1
1	1	1	1	1

The DNF is  $\bar{x}yz + xy\bar{z} + xyz$ 

c)  $F(x, y, z) = x$

$x$	$y$	$z$
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1

The DNF is  $x\bar{y}\bar{z} + x\bar{y}z + xy\bar{z} + xyz$

d)  $F(x, y, z) = x\bar{y}$

$x$	$y$	$z$	$\bar{y}$	$x\bar{y}$
0	0	0	1	0
0	0	1	1	0
0	1	0	0	0
0	1	1	0	0
1	0	0	1	1
1	0	1	1	1
1	1	0	0	0
1	1	1	0	0

The DNF is  $x\bar{y}\bar{z} + x\bar{y}z$ .

**12.2 pg. 822 # 5**

Find the sum-of-products expansion of the Boolean function  $F(w, x, y, z)$  that has the value 1 if and only if an odd number of  $w, x, y,$  and  $z$  have the value 1.

Need to produce all the minterms that have an odd number of 1s. The DNF is simply,

$$\bar{w}xyz + w\bar{x}yz + wx\bar{y}z + wxy\bar{z} + \bar{w}\bar{x}\bar{y}z + \bar{w}\bar{x}y\bar{z} + \bar{w}x\bar{y}\bar{z} + w\bar{x}\bar{y}\bar{z}$$

**12.2 pg. 822 # 11**

Find the product-of-sums expansion of each of the Boolean functions in Exercise 3.

a)  $F(x, y, z) = x + y + z$

Create the table for the function and take the Boolean product of the maxterm where the function evaluates to 0.

$x$	$y$	$z$	$x + y + z$
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

This is already in its CNF,  $x + y + z$ .

b)  $F(x, y, z) = (x + z)y$

$x$	$y$	$z$	$x + z$	$(x + z)y$
0	0	0	0	0
0	0	1	1	0
0	1	0	0	0
0	1	1	1	1
1	0	0	1	0
1	0	1	1	0
1	1	0	1	1
1	1	1	1	1

The CNF is  $(x + y + z)(x + y + \bar{z})(x + \bar{y} + z)(\bar{x} + y + z)(\bar{x} + y + \bar{z})$

c)  $F(x, y, z) = x$

$x$	$y$	$z$
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1

The CNF is  $(x + y + z)(x + y + \bar{z})(x + \bar{y} + z)(x + \bar{y} + \bar{z})$

d)  $F(x, y, z) = x\bar{y}$

$x$	$y$	$z$	$\bar{y}$	$x\bar{y}$
0	0	0	1	0
0	0	1	1	0
0	1	0	0	0
0	1	1	0	0
1	0	0	1	1
1	0	1	1	1
1	1	0	0	0
1	1	1	0	0

The CNF is  $(x + y + z)(x + y + \bar{z})(x + \bar{y} + z)(x + \bar{y} + \bar{z})(\bar{x} + \bar{y} + z)(\bar{x} + \bar{y} + \bar{z})$

### 12.2 pg. 822 # 13

Express each of these Boolean functions using the operators  $+$  and  $\bar{\phantom{x}}$ .

a)  $x + y + z$

$x + y + z$  already satisfies the requirement.

b)  $x + \bar{y}(\bar{x} + z)$

$$\begin{aligned} & x + \bar{y}(\bar{x} + z) \\ &= x + \bar{y}\bar{x} + \bar{y}z \text{ by distributive law} \\ &= x + \overline{(y + x)} + \bar{y}z \text{ by De Morgan's law} \\ &= x + \overline{(y + x)} + \overline{(y + \bar{z})} \text{ by De Morgan's law} \end{aligned}$$

c)  $\overline{x + \bar{y}}$

$\overline{x + \bar{y}}$  already satisfies the requirement.

d)  $\bar{x}(x + \bar{y} + \bar{z})$

$$\begin{aligned} & \bar{x}(x + \bar{y} + \bar{z}) \\ &= (\bar{x}x) + (\bar{x}\bar{y}) + (\bar{x}\bar{z}) \text{ by distributive law} \\ &= 0 + (\bar{x}\bar{y}) + (\bar{x}\bar{z}) \text{ by zero property} \\ &= \overline{(x\bar{y})} + \overline{(x\bar{z})} \text{ by identity law} \\ &= \overline{(x + y)} + \overline{(x + z)} \text{ by De Morgan's law} \\ &= \overline{(x + y)} + \overline{(x + z)} \text{ by De Morgan's law} \end{aligned}$$