

# Control Structures

## ICS312 Machine-Level and Systems Programming

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# Translating high-level structures

- When programming with high-level languages we are used to using high-level structures rather than just branches
  - In fact, many languages don't allow branches (i.e., "goto")
  - C/C++ does though!
- Therefore, it's useful to know how to translate these structures in assembly, so that we can use the same patterns as when writing, say, Java code
  - A compiler does these translations for us with high-level code
  - But in this course, it's not because we write assembly directly that we should ignore everything we've learned with high-level languages and embrace spaghetti code. Quite the opposite.
- Let's start with the most common high-level control structure: if-then-else
  - We already did this in the previous set of slides

# If-then-Else

- A generic if-then-else construct:

```
if (condition) then
    then_block
else
    else_block
```

- Translation into x86 assembly:

```
; instructions to set flags (e.g., cmp ...)
jxx     else_block ; xx so that branch if
                ; condition is false

; code for the then block
jmp endif
else_block:
; code for the else block
endif:
```

# No Else?

- A generic if-then-else construct:

```
if (condition) then
    then_block
```

- Translation into x86 assembly:

```
; instructions to set flags (e.g., cmp ...)
jxx      endif    ; select xx so that branch
              ; if condition is false
; code for the then block
endif:
```

# Else but No Else?

- One can use the pattern in the previous slide even with an else, by doing possibly one extra `mov` instruction

```
if (ebx == 12) {  
    eax = 10;  
} else {  
    eax = 20;  
}
```

```
mov eax, 20  
cmp ebx, 12  
jne endif  
mov eax, 10  
endif:
```

- If `ebx == 12`, the assembly code does an extra `mov` instruction to overwrite the “wrong” value

# Short-circuiting

```
if ((ecx <= 2) || (ebx > 1)) && (edx == 8) {  
    eax = 10;  
} else {  
    eax = 20;  
}
```

- Translating the above to assembly isn't difficult, but some ways of doing it are much easier than others
- A simple way is to try short-circuiting
- For instance, in the above, one can first check whether `edx == 8`, and if not then we know the condition is false
- Let's write the code live right now...

# Short-circuiting

```
if ((ecx <= 2) || (ebx > 1)) && (edx == 8) {  
    eax = 10;  
} else {  
    eax = 20;  
}
```

```
cmp edx, 8  
jne else  
cmp ecx, 2  
jle then  
cmp ebx, 1  
jle else  
then:  
    mov eax, 10  
    jmp end  
else:  
    mov eax, 20  
end:
```

# For Loops

- Let's translate the following loop:

```
sum = 0;
for (i = 0; i <= 10; i++)
    sum += i
```

- Translation

```
    mov eax, 0           ; eax is sum
    mov ebx, 0          ; ebx is i
loop_start:
    cmp ebx, 10         ; compare i and 10
    jg  loop_end        ; if (i>10) go loop_end
    add eax, ebx        ; sum += i
    inc ebx             ; i++
    jmp loop_start      ; goto loop_start
loop_end:
```

# The loop instruction

- It turns out that, for convenience, the x86 assembly provides instructions to do loops!
  - The book lists 3, but we'll talk only about the 1st one
- There is a **loop** instruction
- It is used as: `loop <label>`
- and does
  - Decrement ecx (ecx **has** to be the loop index)
  - If (ecx != 0), branch to the label
    - Only a short jump!
- Let's try to do the loop in our previous example

# For Loops

- Let's translate the following loop:

```
sum = 0;
for (i = 1; i <= 10; i++)
    sum += i
```

- The x86 loop instruction requires that
  - The loop index be stored in ecx
  - The loop index be decremented
  - The loop exits when the loop index is equal to zero
- Given this, we really have to think of this loop in reverse

```
sum = 0
for (i = 10; i > 0; i--)
    sum += i
```

- This loop is equivalent to the previous one, but now it can be directly translated to assembly using the loop instruction

# Using the loop Instruction

- Here is our “reversed” loop

```
sum = 0
for (i = 10; i > 0; i--)
```

sum += i

- And the translation

```
mov  eax, 0           ; eax is sum
mov  ecx, 10          ; ecx is i
loop_start:
add  eax, ecx         ; sum += i
loop loop_start       ; if i > 0 then
                       ; go to loop_start
```

# Using the loop instruction?

- It's totally up to you whether you use the loop instruction or not
  - You should know it though
- In some cases it's very convenient
- In other cases the mental gymnastics needed to modify the loop to use the loop instruction is tedious and error-prone

# While Loops

- A generic while loop

```
while (condition) {  
    body  
}
```

- Translated as:

```
while:  
    ; instructions to set flags (e.g., cmp...)  
    jxx    end_while    ; branches if  
                        ; condition=false  
  
    ; body of loop  
    jmp   while  
end_while:
```

# Do While Loops

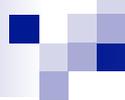
- A generic do while loop

```
do {  
    body  
} while (condition)
```

- Translated as:

do:

```
; body of loop  
; instructions to set flags (e.g., cmp...)  
jxx do ; branches if condition=true
```



# The Compiler

- The compiler does these types of translations for us when we give it high level code
- It can be enlightening to look at the assembly code that the compiler produces
- Let's look at the “Disassembling code with NASM” reading on the course Web site...

# Computing Prime Numbers

- The book has an example of an assembly program that computes prime numbers
- Principle:
  - Try possible prime numbers in increasing order starting at 5
  - Skip even numbers
  - Test whether the possible prime number (the “guess”) is divisible by any number other than 1 and itself
    - If yes, then it’s not a prime, otherwise, it is
- You may want to look at it

# Computing the Sum of an Array

- Let's write a (fragment of a) program that computes the sum of an array
- Let us assume that the array is “declared” in the `.bss` segment as:
  - `array resd 20 ; 20 4-byte values`
- And let us assume that its elements have been set to some values
- We want to compute the numerical sum of all its elements into register `ebx`
- Let's try to write the code together live...

# Computing the Sum of an Array

```
mov    ebx, 0          ; ebx = 0 (sum)
mov    ecx, 0          ; ecx = 0 (loop index)
```

```
main_loop:
```

```
    ; Done?
```

```
    cmp    ecx, 20.    ; compare ecx to 20
```

```
    je     end_main_loop ; if == 20, we're done
```

```
    ; Compute address of current element
```

```
    mov    eax, array  ; eax points to 1st element
```

```
    mov    edx, ecx    ; edx = ecx (loop index)
```

```
    imul  edx, 4       ; edx = 4 * ecx
```

```
    add    eax, edx    ; eax = array + 4 * ecx
```

```
    ; Increment the sum
```

```
    add    ebx, [eax]  ; sum += element
```

```
    ; Move to the next element
```

```
    inc    ecx        ; ecx ++
```

```
    jmp   main_loop   ; loop back
```

```
end_main_loop:
```

# Computing the Sum of an Array

**; SHORTER/SIMPLER VERSION WITH A POINTER**

```
    mov     ebx, 0           ; ebx = 0 (sum)
    mov     ecx, 0           ; ecx = 0 (loop index)
    mov     eax, array       ; eax = array
main_loop:
    ; Done?
    cmp     ecx, 20          ; compare ecx to 20
    je     end_main_loop    ; if == 20, we're done
    ; Increment the sum
    add     ebx, [eax]       ; sum += element
    ; Move to the next element
    add     eax, 4           ; eax += 4
    inc     ecx              ; ecx ++
    jmp    main_loop        ; loop back
end_main_loop:
```

# Going Through a Linked List?

- How about going through a linked list of integers and computing the sum of the integers?
  - Each item in the list contains: (i) `item.value`: a 4-byte integer; (ii) `item.next`: the 4-byte address of the next item (or null)
- Let's try to live code if time...
- The next slide has a version....

# Going Through a Linked List

```
    mov    eax, 0                ; The sum
    mov    ecx, L                ; Our pointer to an item

loop_begin:

    cmp    ecx, 0                ; If nullptr...
    je     end_loop              ; then we're done

    add    eax, [ecx]            ; sum += item.value
    mov    ecx, [ecx + 4]        ; pointer = item.next

    jmp   loop_begin            ; jump back
end_loop:
```

# How about “binning” integers

- (Let's do this if time)
- Say you're given a stream of positive integers between 0 and 299, and the last integer is -1
- Print how many integers are between 0 and 9, how many between 10 and 19, ....., how many are between 290 and 299

# Important Takeaways

- One should avoid spaghetti code in assembly even though we have a “goto”
- High-level control structures are just branch instructions, but using some known/standard patterns
- In an if-then-else one must not forget to “jump over” the second clause
- For for loops, using the `loop` instruction is optional
- While and do-while loops are probably the easiest to implement!

# Conclusion

- Make sure you understand the “sum of an array example” 100%
- Writing control structures in assembly isn't as easy as in high-level languages
- But as long as you follow consistent patterns and use reasonable label names it should be manageable
  
- We have an optional **homework #5**...
- Let's also look at some of the practice problems...
- We'll have an **in-class quiz** on this module next week
- This completes all material for **Midterm #2**