

Subprograms: Arguments

ICS312 Machine-Level and Systems Programming

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Activation Records

- The stack is useful to store and retrieve return addresses, transparently managed via the CALL and RET instructions
- But it's much more useful than this
- In general, when calling a function, one puts all kinds of useful information on the stack
- When the function returns, this information is popped off the stack and the function's caller can safely resume execution
- The set of "useful information" is typically called an **activation record** (or a "**stack frame**")
- One very important component of an activation record is the **parameters** passed to the function
 - Another is the **return address**, as we've already seen

Subprogram Conventions

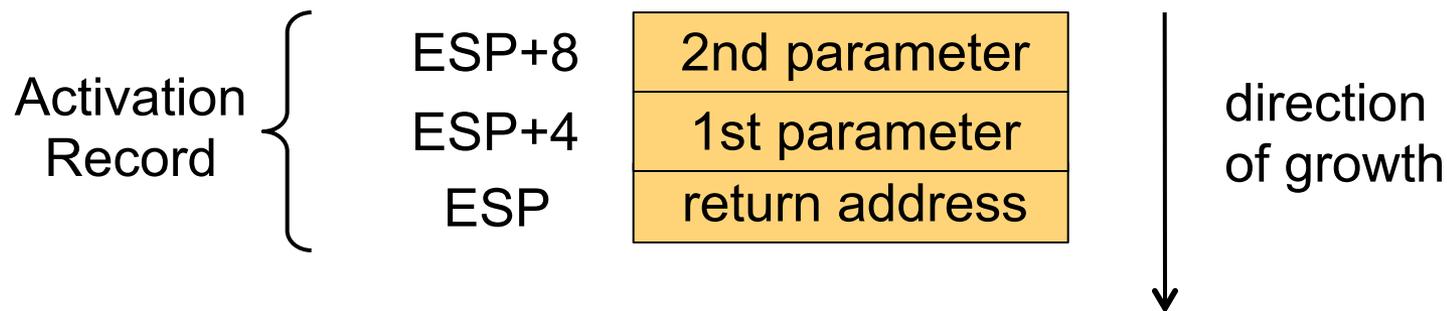
- When writing assembly, you could do whatever you want
- For instance, you could devise a clever scheme that reuses register values in creative ways instead of the stack
- Such solutions are typically error prone, making the code difficult to debug/extend/maintain, but can enhance performance
- Typically, one uses a consistent **calling convention**, so that there is a generic way to call a subprogram
- Of course compilers use calling conventions
 - The compiler, when generating assembly code, follows a standard method to generate assembly for all function calls
- Some languages specify which calling convention should be used
- **What we describe in all that follows is (mostly) the convention used by the C language**
 - i.e., C compilers must use this convention when generating assembly code from C code
 - We'll also use this convention when writing assembly by hand

A Simple Activation Record

- To call a function you have to follow these steps:
 - **Push the parameters onto the stack**
 - Execute the CALL instruction, which pushes the return address onto the stack
- Warning: In the C calling convention parameters are pushed onto the stack **in reverse order!**
 - Say the function is $f(a,b,c)$
 - c is pushed onto the stack first
 - b is pushed onto the stack second
 - a is pushed onto the stack third
 - Makes sense-ish: the first parameter should be at a lower address than the second parameter

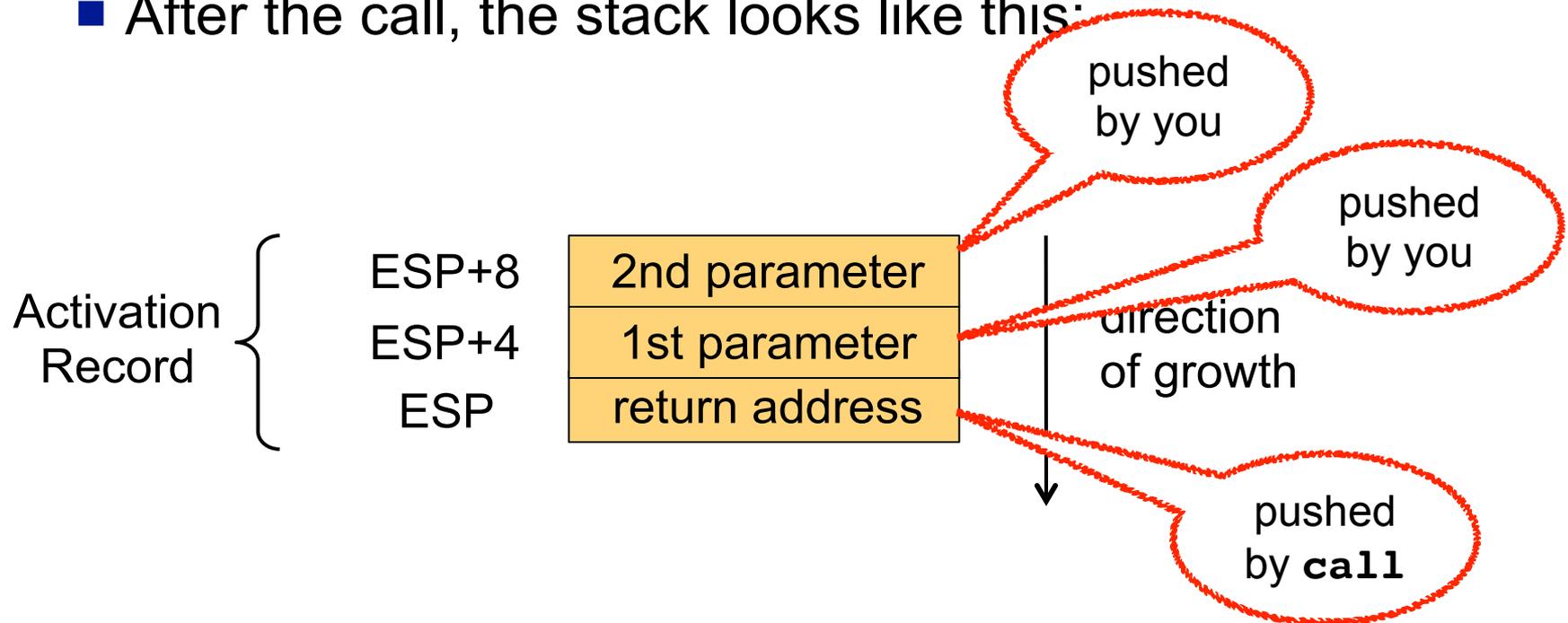
A Simple Activation Record

- Say you want to **call** a function with two 32-bit parameters
 - If parameters are < 32 bits, they need to be extended to 32-bit values (movzx or movsx), at least in this course
- After the call, the stack looks like this:



A Simple Activation Record

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Using the Parameters

- Inside the code of the subprogram, parameters can be accessed via indirection from the stack pointer
- In our previous example:
 - `mov eax, [ESP + 4]` ; puts 1st parameter into eax
 - `mov ebx, [ESP + 8]` ; puts 2nd parameter into ebx
- Typically the subprogram does not pop the parameters off the stack before using them
 - It would be annoying to have to pop the return address first, and then push it back
 - It's convenient to have the parameters always stored in memory as opposed to being careful to constantly preserve them in registers
 - They may be copied into registers for performance reasons
 - But we can always get their original values from the stack

Accessing the stack in C

```
void main(int x) {  
    x++;    // translated as: inc dword [esp + 4]  
}
```

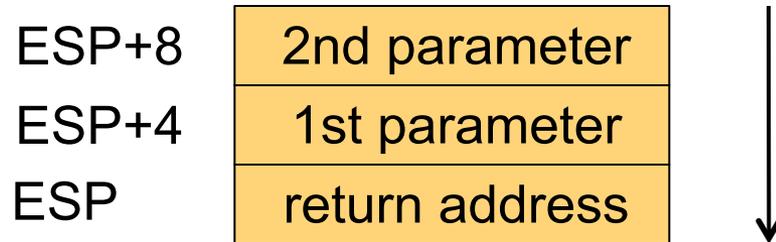
- The activation record on the stack is the subprogram's little play pen
 - And yes, you can add one to a parameter as seen above, just as if it were a local variable
- The subprogram can modify its activation record in many ways, and eventually its activation record is wiped out anyway
- But, turns out, there is still a problem...

ESP and EBP

- There is one problem with referencing parameters using ESP, as in [ESP+4]
- If the subprogram uses the stack for something else, ESP will be modified!
 - So at some point in the program, the 1st parameter should be accessed as [ESP+4]
 - And at some other point, it may be accessed as [ESP+8], [ESP+12], etc., depending on how the stack grows
- So the convention is to use the **EBP** register as an anchor to save the value of ESP as soon as the subprogram starts
- Afterwards, the 1st parameter is **always** accessed as [EBP+4], the 2nd parameter is **always** accessed as [EBP+8], etc.

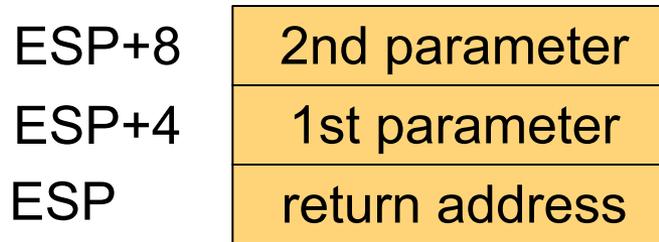
ESP and EBP

- Stack as it is when the subprogram begins

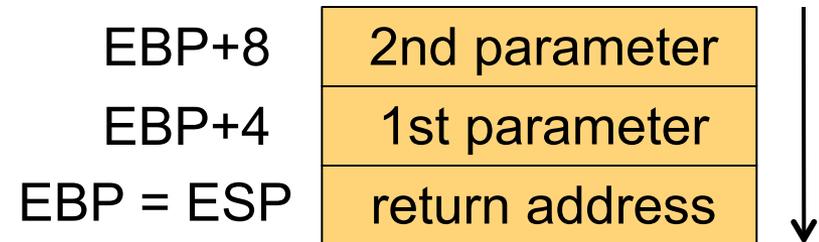


ESP and EBP

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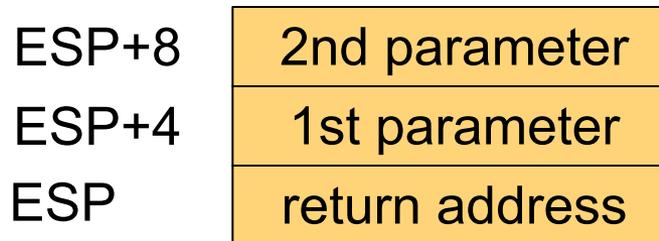


- EBP = ESP

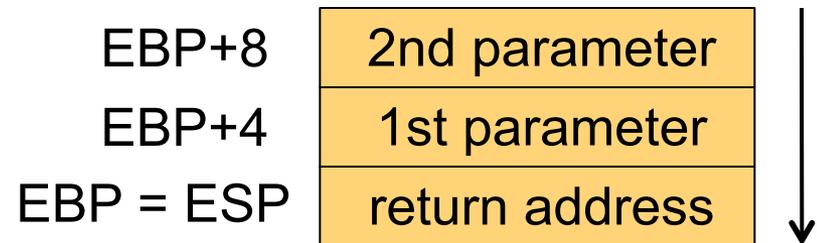


ESP and EBP

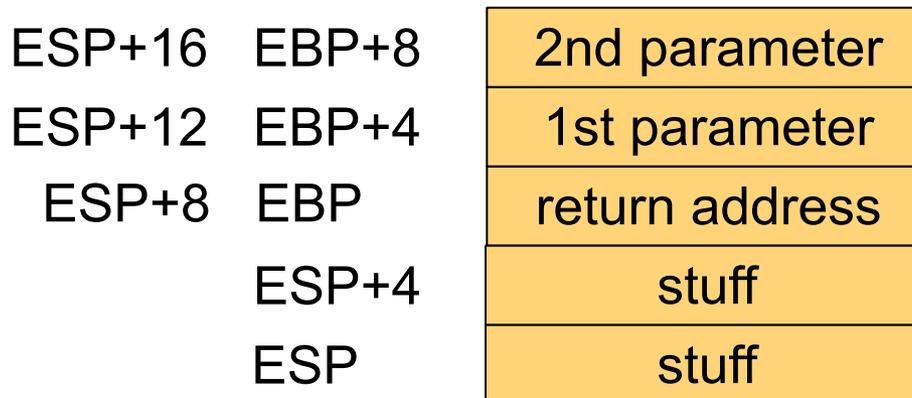
- Stack as it is when the subprogram begins



- EBP = ESP



- Further use of the stack



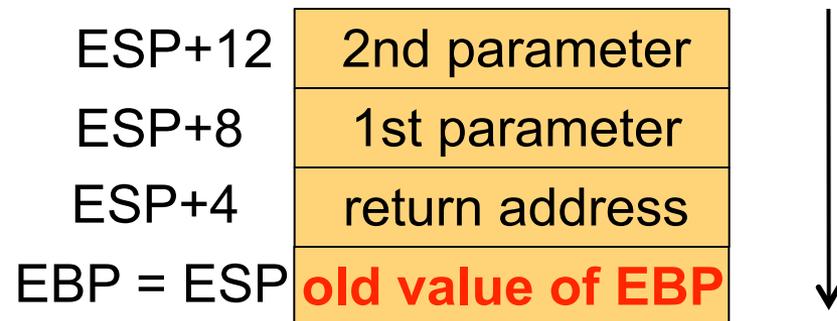
Parameters still referred to as EBP+4 and EBP+8

ESP and EBP Mayhem

- **Big problem:** The **caller** may have been using EBP!
 - Typically to access its own parameters!!!
- So you can't just overwrite EBP (you = a subprogram being called)
- Because when you return, the caller will have a wrong EBP and will access its own parameters erroneously
- So you have to save EBP before overwriting it for your own purposes
- How do we deal with having to save stuff?
- We use the stack!!

Saving EBP on the Stack

- The convention is to first **push the value of EBP** onto the stack and then set **EBP = ESP**, as soon as the program starts
- So, the stack before the subprogram truly begins is:



- Parameter accesses:
 - 1st parameter: [EBP+8]
 - 2nd parameter: [EBP+12]

- At the end of the subprogram, the value of EBP is popped and restored with a simple POP instruction

Subprogram Skeleton

`func:`

```
push    ebp        ; save my caller's EBP
```

```
mov     ebp, esp   ; set EBP = ESP
```

```
. . .                ; subprogram code
```

```
pop     ebp        ; restore my caller's EBP
```

```
ret     ; returns
```

Returning from a Subprogram

- After the subprogram finishes, one must “clean up” the stack
- The stack has on it:
 - The old EBP value, the return address, the parameters
- The old EBP value is popped in the subprogram (at the end)
- The return address is removed by the RET instruction
 - You don't see the POP, but it's there
- **The parameters need to be removed from the stack**
- The C convention specifies that the **caller code must remove the parameters from the stack**
 - Other languages specify that the callee must do it
 - In fact, it is well-known that it's a little bit more efficient to have the subprogram (i.e., the callee) do it!
- So one may wonder why C opts for the slower approach
- Turns out, it's all because of *varargs*
 - *Let's go into a bit of a detour.... if you're confused already, you can safely skip the next 2 slides when you study this content*

Variable Number of Arguments

- C allows or the declaration of functions with variable number of arguments
- A well-known example: `printf()`
 - `printf("%d", 2);`
 - `printf("%d %d", 2, 3);`
 - `printf("%s %d %c %f", "foo", 1, 'f', 3.14);`
- So sometimes there will be 1 argument to remove from the stack, sometimes 2, sometimes 3, etc.
- Having the subprogram (in this case `printf`) remove the arguments from the stack requires some complexity
 - e.g., pass an extra (shadow) parameter that specifies how many arguments should be removed
- Instead, the convention is that the caller removes the arguments, because it always knows how many there are
 - e.g., it's easy for a compiler to generate code that does this

Variable of Arguments in C

- Just in case you are curious, here is an example of a C program with a vararg function

```
#include <stdarg.h>
#include <stdio.h>

int func(int first, ...) {
    va_list args;
    va_start(args, first);
    printf("arg #1 = %d\n", first);
    printf("arg #2 = %d\n", va_arg(args, int));
    printf("arg #3 = %s\n", va_arg(args, char*));
    va_end(args);
}
```

```
int main() {
    func(2, (void*)3,
        (void*)"foo");
}
```

Vararg functions are a bit dangerous. If you call `va_arg()` more times than there are arguments on the stack, you'll just get bogus values!

Example: Calling a Subprogram

Caller:

```
push    dword    2        ; second parameter
push    dword    1        ; first parameter
call    func          ; call the function
add    esp, 8        ; pop the two arguments
```

- Note that to pop the two arguments we merely add 8 to the stack pointer ESP
 - Since we do not care to get the values of the arguments at this point, it's quicker than to call pop twice!
 - This is one case in which we do modify ESP directly
- This is because the stack is not really a stack, but just an array with a pointer: the parameter values will be overwritten next time a function is called or next time the stack is used
 - We don't zero out "old" value, we just lazily overwrite them later

Return Values?

- Often, one wants a subprogram to return a value
 - e.g., a function that computes some number
- There are several ways to do this
- One way is to pass as a parameter the address of a zone of memory in which some result should be written
 - As in: `void foo(int *x);` `foo(&a);`
- This is not a *true* return value
 - As in: `int foo();`
- **The C convention is that the return value is always stored in EAX when the function returns**
 - It's the responsibility of the caller to save the EAX value before the call (if needed) and to restore it later

Recall the NASM Skeleton

```
    ; include directives

segment .data
    ; DX directives

segment .bss
    ; RESX directives

segment .text
    global asm_main
asm_main:
    enter    0,0
    pusha
    ; Your program here
    popa
    mov     eax, 0
    leave
    ret
```

Returns value 0



Recall the NASM Skeleton

```
    ; include directives

segment .data
    ; DX directives

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    ; RESX directives

segment .text
    global asm_main
asm_main:
    enter 0,0
    pusha
    ; Your program here
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    mov  eax, 0
    leave
    ret
```

The last two remaining things that we haven't explained yet (but soon)

A Full Example

```
L      dd  42, 43, 44, 45, 56
...
push   dword L
call   func
add    esp, 4
call   print_int
...
func:
    push   ebp
    mov    ebp, esp
    push   dword [ebp+8]
    push   8
    call   reference
    add    esp, 8
    add    eax, 10
    pop    ebp
    ret

reference:
    push   ebp
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    mov    eax, [eax]
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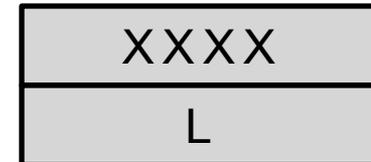


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ESP →



A Full Example

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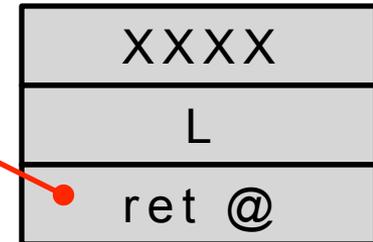
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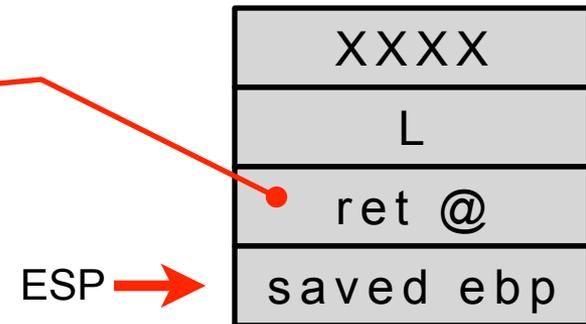
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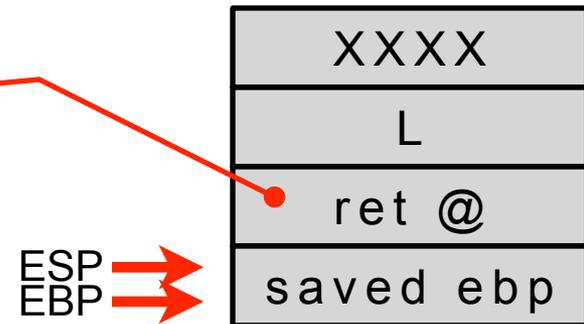
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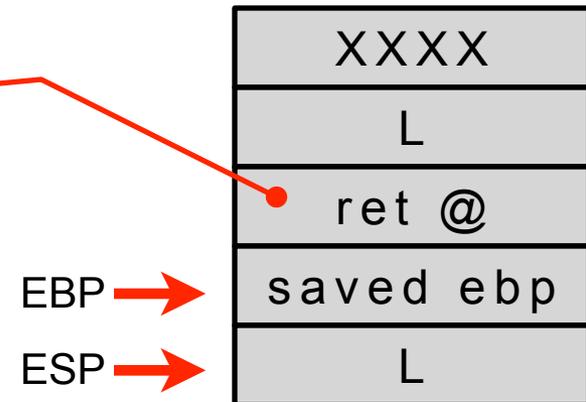
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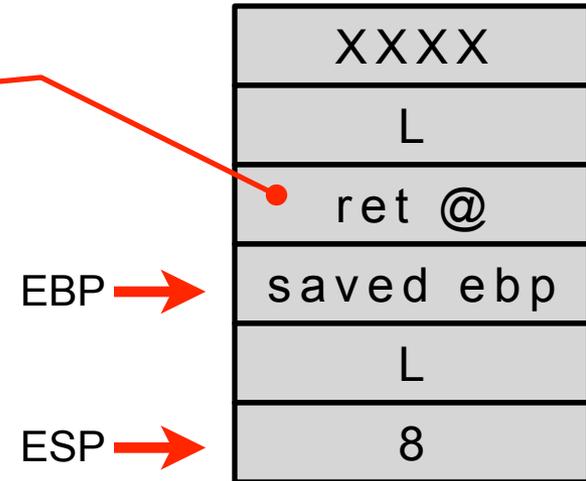
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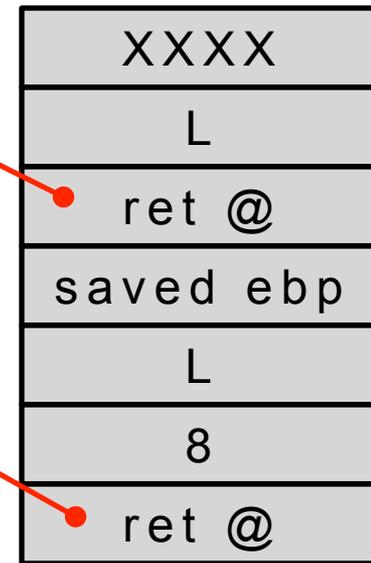
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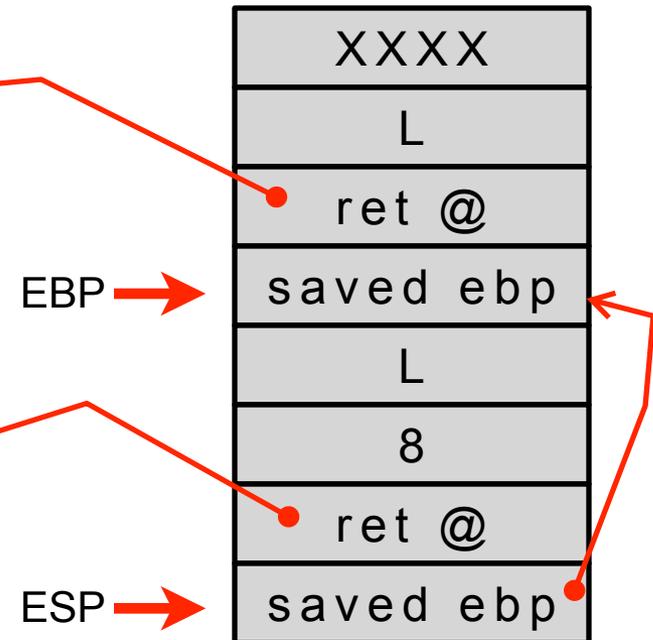
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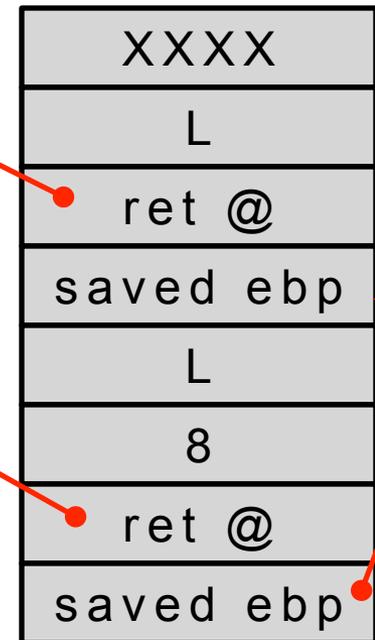
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func:

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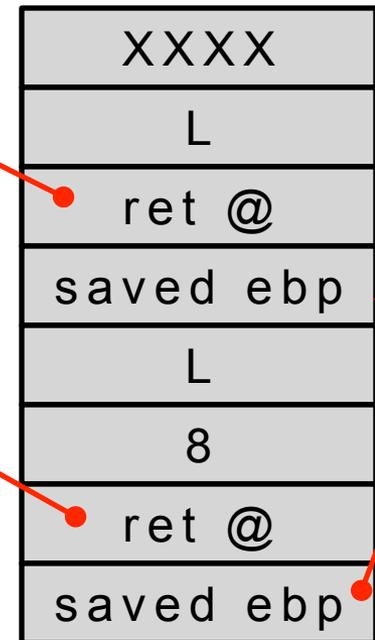
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```



ESP →
EBP →

EAX = L

A Full Example

```
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```

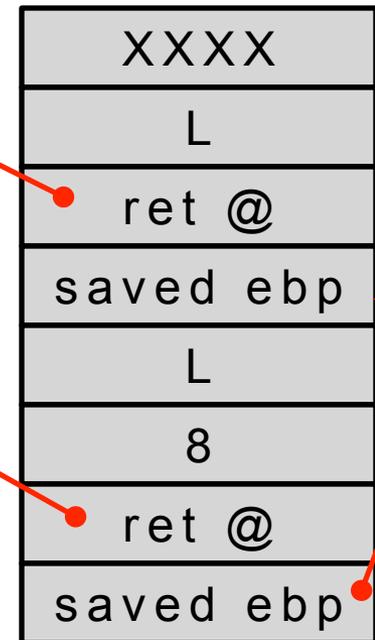
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func:

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add    eax, [ebp+8]  
mov    eax, [eax]  
pop    ebp  
ret
```



EAX = L + 8

A Full Example

```
L      dd  42, 43, 44, 45, 56
```

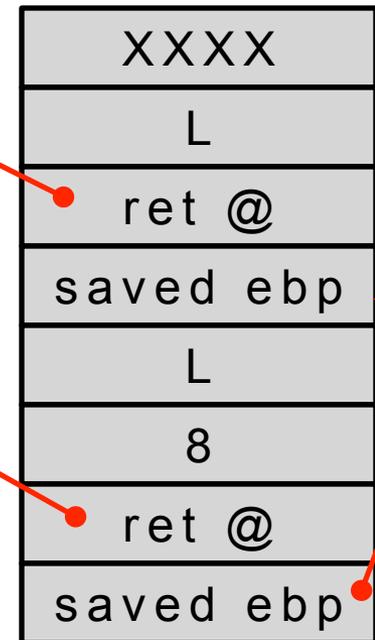
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...  
push   dword L  
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func:
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ret
```



ESP →
EBP →

$$\text{EAX} = [\text{L} + 8] = 44$$

A Full Example

```
L      dd  42, 43, 44, 45, 56
```

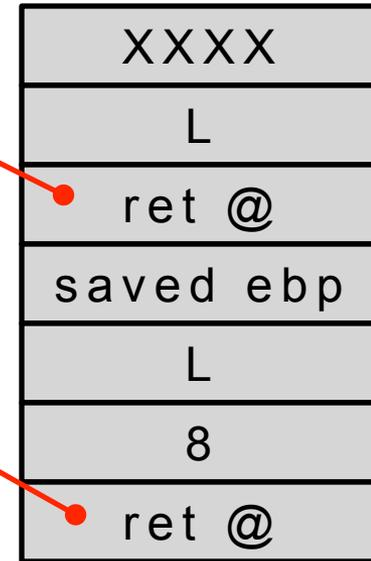
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```



EAX = 44

A Full Example

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```

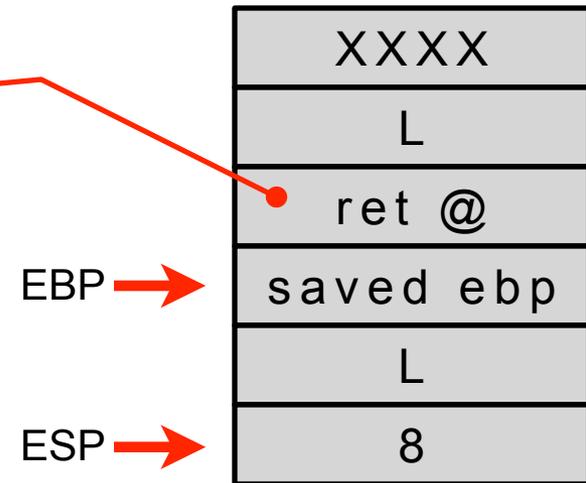
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push   ebp  
mov    ebp, esp  
mov    eax, [ebp+12]  
add    eax, [ebp+8]  
mov    eax, [eax]  
pop    ebp  
ret
```



EAX = 44

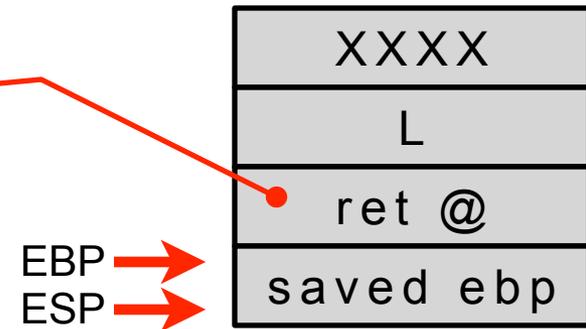
A Full Example

```
L      dd  42, 43, 44, 45, 56
...
push   dword L
call   func
add    esp, 4
call   print_int
...
func:
```

```
    push   ebp
    mov    ebp, esp
    push   dword [ebp+8]
    push   8
    call   reference
    add    esp, 8
    add    eax, 10
    pop    ebp
    ret
```

reference:

```
    push   ebp
    mov    ebp, esp
    mov    eax, [ebp+12]
    add    eax, [ebp+8]
    mov    eax, [eax]
    pop    ebp
    ret
```



EAX = 44

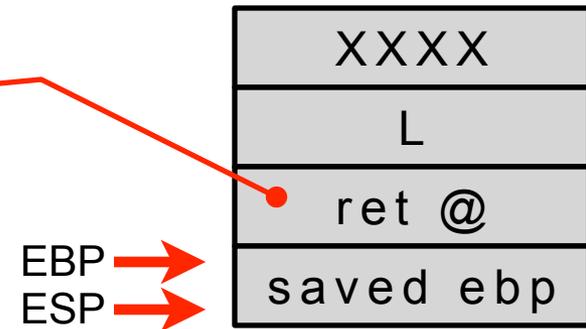
A Full Example

```
L      dd  42, 43, 44, 45, 56
...
push   dword L
call   func
add    esp, 4
call   print_int
...
func:
```

```
    push   ebp
    mov    ebp, esp
    push   dword [ebp+8]
    push   8
    call   reference
    add    esp, 8
    add    eax, 10
    pop    ebp
    ret
```

reference:

```
    push   ebp
    mov    ebp, esp
    mov    eax, [ebp+12]
    add    eax, [ebp+8]
    mov    eax, [eax]
    pop    ebp
    ret
```



$$EAX = 44 + 10 = 54$$

A Full Example

```
L      dd  42, 43, 44, 45, 56
```

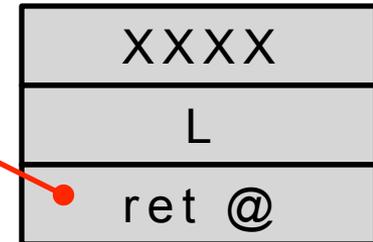
```
...  
push   dword L  
call   func  
add    esp, 4  
call   print_int  
...
```

```
func:
```

```
push   ebp  
mov    ebp, esp  
push   dword [ebp+8]  
push   8  
call   reference  
add    esp, 8  
add    eax, 10  
pop    ebp  
ret
```

```
reference:
```

```
push   ebp  
mov    ebp, esp  
mov    eax, [ebp+12]  
add    eax, [ebp+8]  
mov    eax, [eax]  
pop    ebp  
ret
```



EAX = 54

A Full Example

```
L      dd  42, 43, 44, 45, 56
```

```
...  
push   dword L  
call   func  
add    esp, 4  
call   print_int  
...
```

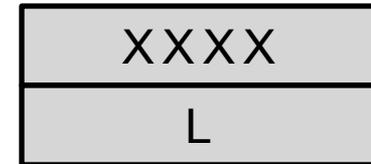
```
func:
```

```
    push   ebp  
    mov    ebp, esp  
    push   dword [ebp+8]  
    push   8  
    call   reference  
    add    esp, 8  
    add    eax, 10  
    pop    ebp  
    ret
```

```
reference:
```

```
    push   ebp  
    mov    ebp, esp  
    mov    eax, [ebp+12]  
    add    eax, [ebp+8]  
    mov    eax, [eax]  
    pop    ebp  
    ret
```

ESP →



EAX = 54

A Full Example

```
L      dd  42, 43, 44, 45, 56
```

```
...
```

```
push   dword L
```

```
call   func
```

```
add    esp, 4
```

```
call   print_int
```

```
...
```

```
func:
```

```
push   ebp
```

```
mov    ebp, esp
```

```
push   dword [ebp+8]
```

```
push   8
```

```
call   reference
```

```
add    esp, 8
```

```
add    eax, 10
```

```
pop    ebp
```

```
ret
```

```
reference:
```

```
push   ebp
```

```
mov    ebp, esp
```

```
mov    eax, [ebp+12]
```

```
add    eax, [ebp+8]
```

```
mov    eax, [eax]
```

```
pop    ebp
```

```
ret
```

ESP →

XXXX

EAX = 54

A Full Example

```
L      dd  42, 43, 44, 45, 56
...
push   dword L
call   func
add    esp, 4
call   print_int
...
func:
    push   ebp
    mov    ebp, esp
    push   dword [ebp+8]
    push   8
    call   reference
    add    esp, 8
    add    eax, 10
    pop    ebp
    ret

reference:
    push   ebp
    mov    ebp, esp
    mov    eax, [ebp+12]
    add    eax, [ebp+8]
    mov    eax, [eax]
    pop    ebp
    ret
```



prints "54"

C Translation of the previous program (reverse-engineering)

```
#include <stdio.h>

int    L[5] = {42, 43, 44, 45, 56};
int func(int *array);
int reference(int a, int *ptr);

int main(int argc, char **argv) {
    // ...
    printf("%d", func(L));
    // ...
}

int func(int *array) {
    return 10 + reference(8, array); // The 8 should be a 2 (2nd element of the array)
}

int reference(int a, int *ptr) {
    return *((int *)((char *)ptr + a));
    // return ptr[a / sizeof(int)];
}
```

In-class Exercise

- What 3 (perhaps 4) things are wrong with the following program?

```
push    ebx
push    dword 30
call    func
add     esp, 4
call    print_int
call    print_nl
. . .
```

```
func:   push    ebp
        mov     ebp, esp
        mov     eax, [ebp+8]
        add     eax, [ebp+4]
        ret
```

In-class Exercise

- What 4 things are wrong with the following program?

```
push    ebx
push    dword 30
call    func
add     esp, 8
call    print_int
call    print_nl
. . .
```

```
func:   push    ebp
        mov     ebp, esp
        mov     eax, [ebp+12]
        add     eax, [ebp+8]
        pop     ebp
        ret
```

In-class Exercise

- What does the stack look like?

```
push    ebx
push    dword 30
call    func
        <----- THERE?
add     esp, 8
call    print_int
call    print_nl
. . .
```

```
func:   push    ebp
        mov     ebp, esp
        <----- HERE?
        mov     eax, [ebp+12]
        add     eax, [ebp+8]
        pop     ebp
        ret
```

In-class Exercise

- What does the stack look like?

```
push    ebx
push    dword 30
call    func
```

xxxxxx
EBX
30

```
    <-----
add     esp, 8
call    print_int
call    print_nl
. . .
```

```
func:  push    ebp
       mov     ebp, esp
```

xxxxxx
EBX
30
Return @
EBP

```
    <-----
mov     eax, [ebp+12]
add     eax, [ebp+8]
pop     ebp
ret
```

A Full Example with Subprograms

- The book has a full example in Section 4.5.1
- Let's do another example here
- Say we want to write a program that first reads in a sequence of 10 integers and then prints the number of odd integers
- We will use three functions:
 - `get_integers()`: get the 10 integers from the user
 - `count_odds()`: count the number of odd integers
 - `is_odd()`: determines whether an integer is odd
- We could do this without functions
 - The code would most likely be less readable
 - But faster! (usual tradeoff)
- For now, we're writing the code in the most modular and "clean" fashion
- Let's first look at the easy main program

Example: Main program

```
%include "asm_io.inc"
```

```
segment .data
```

```
    msg_odd db    "The number of odd numbers is: ",0
```

```
segment .bss
```

```
    integers resd 10 ; space for 10 integers
```

```
segment .text
```

```
    global asm_main
```

```
asm_main:
```

```
    enter    0,0      ; set up
```

```
    pusha                ; set up
```

```
    popa                ; clean up
```

```
    mov     eax, 0      ; clean up
```

```
    leave               ; clean up
```

```
    ret                ; clean up
```

```
    push    integers    ; we pass integers (address) to get_integers
    push    dword 10     ; we pass the number of integers to get_integers
    call   get_integers  ; call get_integers
    add    esp, 8        ; clean up the stack
    mov    eax, msg_odd  ; store the address of the message to print into eax
    call   print_string  ; print the message
    push    integers    ; we pass integers (address) to count_odds
    push    dword 10     ; we pass the number of integers to count_odds
    call   count_odds    ; call count_odds
    add    esp, 8        ; clean up the stack
    call   print_int     ; print the content of eax as an integer
                                ; (this is what count_odds returned)
    call   print_nl     ; print a new line
```

Piecemeal segment declarations

- The NASM assembler allows for the declaration of multiple `.data`, `.bss`, and `.text` segments
- This makes it possible to declare subprograms in their own region of the `.asm` file, with parts of `.data` and `.bss` segments that are relevant for the subprograms
- Let's look at the `get_integers()` subprogram

Example: get_integers

```
; FUNCTION: Get_Integers
; Takes two parameters: an address in memory in which to store integers, and a number of integers to store (>0)
; Destroys values of eax, ebx, and ecx!!
```

```
segment .data
```

```
    msg_int    db    "Enter an integer: ",0
```

```
segment .text
```

```
get_integers:
```

```
    push  ebp            ; save the value of EBP of the caller
    mov   ebp, esp      ; update the value of EBP for this subprogram
```

```
    mov   ecx, [ebp + 12] ; ECX = address at which to store the integers (parameter #2)
    mov   ebx, [ebp + 8]  ; EBX = number of integers to read (parameter #1)
    shl   ebx, 2          ; EBX = EBX * 4 (unsigned)
    add   ebx, ecx        ; EBX = ECX + EBX = address beyond that of the last integer to be stored
```

```
loop1:
```

```
    mov   eax, msg_int   ; EAX = address of the message to print
    call  print_string;  ; print the message
    call  read_int       ; read an integer from the keyboard (which will be stored in EAX)
    mov   [ecx], eax     ; store the integer in memory at the correct address
    add   ecx, 4         ; ECX = ECX + 4
    cmp   ecx, ebx      ; compare ECX, EBX
    jb   loop1          ; if ECX < EBX, jump to loop1 (unsigned)
```

```
    pop   ebp           ; restore the value of EBP
    ret                ; clean up
```

Example: count_odds

```
; FUNCTION: count_odds
; Takes two parameters: an address in memory in which integers are stored, and the number of integers (>0)
; Destroys values of eax, ebx, and edx!! (eax = returned value)
```

```
segment .text
```

```
count_odds:
```

```
    push    ebp                ; save the value of EBP of the caller
```

```
    mov     ebp, esp          ; update the value of EBP for this subprogram
```

```
    mov     eax, [ebp + 12]    ; EAX = address at which integers are stored (parameter #2)
```

```
    mov     ebx, [ebp + 8]     ; EBX = number of integers (parameter #1)
```

```
    shl     ebx, 2            ; EBX = EBX * 4 (unsigned)
```

```
    add     ebx, eax          ; EBX = EAX + EBX = address beyond that of the last integer
```

```
    sub     ebx, 4            ; EBX = EBX - 4 = address of the last integer
```

```
    xor     edx, edx          ; EDX = 0 = number of odd integers
```

```
loop2:
```

```
    push   dword [ebx]        ; store the current integer on the stack
```

```
    call  is_odd              ; call is_odd
```

```
    add   esp, 4              ; clean up the stack
```

```
    add   edx, eax            ; EDX += EAX (EAX = 0 if even, EAX = 1 if odd)
```

```
    sub   ebx, 4              ; EBX = EBX - 4
```

```
    cmp   ebx, [ebp+12]      ; compare EBX and the address of the first integer
```

```
    jnb  loop2                ; if EBX >= [EBP+12] jump to loop2 (unsigned test)
```

```
    mov   eax, edx            ; EAX = EDX (= number of odd integers)
```

```
    pop   ebp                ; restore the value of EBP
```

```
    ret                                ; clean up
```

Example: is_odd

```
; FUNCTION: is_odd  
; Takes one parameter: an integers (>0)  
; Destroys values of eax and ecx (eax = returned value)
```

```
segment .text
```

```
is_odd:
```

```
    push    ebp                ; save the value of EBP of the caller  
    mov     ebp, esp          ; update the value of EBP for this subprogram  
  
    mov     eax, 0            ; EAX = 0  
    mov     ecx, [ebp+8]      ; ECX = integer (parameter #1)  
    shr     ecx, 1            ; Right logical shift  
    adc     eax, 0            ; EAX = EAX + carry (if even: EAX = 0, if odd: EAX = 1)  
  
    pop     ebp              ; restore the value of EBP  
    ret
```

```
; clean up
```

Destroyed Registers?

- Note that in the previous program we have added comments specifying which registers are destroyed
- The caller is then responsible for making sure that its registers are not corrupted
- However, in a program that has many functions it becomes really annoying to constantly have to pay attention to what needs to be saved and what doesn't
- **The typical approach is to have the subprogram save what it knows it will overwrite onto the stack!**
 - And comment that the caller doesn't need to worry about anything
- Let's look at examples...

Saving Registers in Subprograms

- Just saving EBP

func:

```
push    ebp        ; save original EBP
mov     ebp, esp   ; set EBP = ESP

. . .           ; subprogram code

mov     eax, ...   ; set return value

pop     ebp        ; restore original EBP
ret      ; returns
```

Saving Registers in Subprograms

- Saving, for instance, EBX and ECX, in addition to EBP

func:

```
push    ebp           ; save original EBP
mov     ebp, esp     ; set EBP = ESP
push    ebx           ; save EBX
push    ecx           ; save ECX

. . .                ; subprogram code

mov     eax, ...     ; set return value

pop     ecx           ; restore ECX
pop     ebx           ; restore EBX
pop     ebp           ; restore ebp
ret                               ; returns
```

Saving Registers in Subprograms

- Saving “all” registers using PUSHAD and POPAD

func:

```
push    ebp           ; save original EBP
mov     ebp, esp     ; set EBP = ESP
pushad  ; save all (including new EBP)

. . .                ; subprogram code

mov     eax, ...     ; set return value

popad   ; restore all (including new EBP)
pop     ebp          ; restore original ebp
ret     ; returns
```

Saving Registers in Subprograms

- Saving “all” registers using PUSHA and POPA

func:

```
push    ebp           ; save original EBP
mov     ebp, esp      ; set EBP = ESP
pusha   ; save all (including new EBP)

. . .                ; subprogram code

mov     eax, ...      ; set return value

popa    ; restore all (including new EBP)
pop     ebp           ; restore original ebp
ret     ; returns
```

Overwrites the return value
that's stored in eax!

Dealing with Return Value

- Saving “all” registers using PUSHAD and POPAD + return value handling

```
.bss:
    returnvalue    resd    1        ; place in memory for the return value
func:
    push          ebp             ; save original EBP
    mov           ebp, esp       ; set EBP = ESP
    pushad        ; save all (including new EBP)

    . . .                        ; subprogram code

    mov           [returnvalue], eax ; save return value in memory

    popad         ; restore all (including new EBP)
    mov           eax, [returnvalue] ; retrieve the saved return
value
                                ; (as done in our skeleton)

    pop          ebp             ; restore original ebp
    ret
```

Dealing with Return Value

- Saving “all” registers using PUSHAD and POPAD + return value handling

```
.bss:
    returnvalue    resd    1        ; place in memory for the return value
func:
    push    ebp
    mov     ebp, esp
    pushad
    . . .
    mov     [returnvalue], eax      ; store return value in memory

    popad
    mov     eax, [returnvalue]      ; retrieve the saved return value
    ; (as done in our skeleton)

    pop     ebp                    ; restore original ebp
    ret
```

A much better option is to put the return value in a **local variable**, which we'll see in the next set of lecture notes

Recursion

- The subprogram calling conventions we have just described enable recursion out of the box!
- Let's live-code an example program that computes the sum of the first n integers
 - Yes, it's $n(n+1)/2$, and even if we didn't know this, an iterative program would be more efficient; but for the sake of this example let's just write a recursive program to compute it

Example: Recursive Program

```
. . .
segment .data
    msg1        db        `Enter n: `, 0
    msg2        db        `The sum is: `, 0
segment .text
. . .                ; declaration of asm_main and setup

mov                eax, msg1        ; eax = address of msg1
call               print_string     ; print msg1
call               read_int         ; get an integer from the keyboard (in EAX)
push               eax              ; put the integer on the stack (parameter
#1)

call               recursive_sum     ; call recursive_sum
add                esp, 4           ; remove the parameter from the stack
mov                ebx, eax         ; save the value returned by recursive_sum
mov                eax, msg2        ; eax = address of msg2
call               print_string     ; print msg2
mov                eax, ebx         ; eax = sum
call               print_int        ; print the sum
call               print_nl         ; print a new line

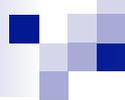
. . .                ; cleanup
```

Example: recursive_sum()

```
segment .bss
    value            resd, 1            ; to store the return value temporarily
segment .text
recursive_sum
    push            ebp                ; save ebp
    mov             ebp, esp          ; set EBP = ESP
    pusha           ; save all registers (probably overkill)
    mov             ebx, [ebp+8]      ; ebx = integer (parameter #1)
    cmp             ebx, 0            ; ebx = 0 ?
    jnz            next              ; if (ebx != 0) go to next
    xor             ecx, ecx          ; ECX = 0
    jmp            end                ; Jump to end
next:
    mov             ecx, ebx          ; ECX = EBX
    dec             ecx              ; ECX = ECX - 1
    push            ecx              ; put ECX on the stack
    call            recursive_sum     ; recursive call to recursive_sum!
    add             esp, 4            ; pop the parameter from the stack
    add             ebx, eax          ; EBX = EBX + recursive_sum(EBX -1)
    mov             ecx, ebx          ; ECX = EBX
end:
    mov             [value], ecx      ; at this point, ECX contains the result
    popa            ; restore registers
    mov             eax, [value]      ; put the saved returned value into eax
    pop             ebp              ; restore EBP
    ret                ; return
```

Important Takeaways

- A subprogram's data is on the stack and called an activation record or a stack frame
 - So far we have seen this content in the activation records: return address, saved EBP value, other saved register values as needed, arguments
- Arguments are
 - Pushed by the caller (in “reverse order” from a high-level perspective)
 - Popped by the caller
- The EBP register is used as an “anchor” so that the code of the subprogram always knows where its activation record is
 - Needed, because the value of ESP keeps changing
- The EBP value of a subprogram's caller should be saved on the stack by that subprogram before it overwrites that value
- Values of other registers used by the caller may also need to be pushed onto the stack
- A subprogram's return value is store in the EAX register
- And just like that, we have recursion “for free”



Conclusion

- You must absolutely make sure you fully understand all code examples in this set of slides
 - Not that this is not true for all code examples in this course!
- In the next set of lecture notes we'll talk about local variables in subprograms