## **Midterm Review**

# ICS332 Operating Systems

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## What to Expect

- Closed notes/computer/phone
- Scope: All modules up to and including Synchronization
- Material to review:
  - Quiz solutions
  - Homework solutions
  - Lecture notes
  - Reading assignments in the textbook
    - Especially the examples

## **What Questions to Expect**

- Some quiz-like
- Some short "how?" or "why?" questions
  - Answer in a few sentences
- Study for the above by going through the material pretending you're a professor who has to come up with a bunch of quiz / short questions
- Some may have to do with what you did in the homework assignments
- They will be no "write a program" question
  - □ At most 1-2 lines of code to insert/fix in a given program
- Questions like the homework assignments
- There should be very few surprises



### **Example Short Question**

What's the difference between an Interrupt and a Trap? Give an example of each

- What's the difference between an Interrupt and a Trap? Give an example of each
  - An interrupt is dan external event
    - Example: keyboard input, disk operation completion, mouse click, network packet arrival, etc.
  - A trap is an internal event caused by an instruction
    - Example: divide by zero, illegal memory access, system call instruction



- What memory reads/writes happen when a context-switch occurs?
  - Say we context switch from process A to process
  - A's registers are written into its PCB
  - B's registers are read from its PCB

- What's a Zombie?
- What's an Orphan?

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- What's a Zombie?
  - A process that has terminated (i.e., called exit()), but whose parent has not acknowledged the death (i.e., not call to wait(), waitpid()).
  - It is kept around so that later the parent can find out its return value
- What's an Orphan?
  - A process whose parent has died
  - In Linux it is adopted by process PID=1 (meaning that its PPID=1)

- What is priority inversion?
- Why should context-switching overhead be low?
- What happens if at the end of a time quantum of a running process the ready queue is empty?
- A large time quantum value in a round-robin scheduler will be preferred by what kind of processes?
- And so on....

## **Longer Questions**

- Longer questions will ressemble homework and in-class exercises / examples
  - Given a program with calls to fork(), exec(), wait(), dup(), close(), etc. what is wrong / what does it do?
  - Given a schedule what can you say about the scheduling algorithm?
  - Given a set of jobs with CPU and I/O burst times, what do various scheduling algorithms do?
  - Given a multi-threaded programs, what does it do? What's wrong with it?
    - How would you add locks?
    - Is there a deadlock?

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## **Synchronization Questions**

- We did not have a homework in the Synchronization module
- So let's review this right now
- Main concepts
  - Race conditions
  - Locks
    - Spinning vs. Blocking
  - Deadlocks

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#### **Race Condition**

- Two threads, one global variable a = 0;
- Thread #1 does: a+=2;
- Thread #2 does: a-=1;
- What are the possible final values of a?

#### **Race Condition**

- Two threads, one global variable a = 0;
- Thread #1 does: a+=2;
- Thread #2 does: a-=1;
- What are the possible final values of a?
- The "clean" execution is a = 1
- Each thread could have a lost update, which would lead to -1 or 2
- Answer: {-1, 1, 2}

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#### Locks

Consider the following code fragment:

```
1 int a = 0;
2
3 void f(int x) {
4    x++;
5    a = a + x;
6 }
```

Where would you add lock() and unlock() calls so that multiple threads can safely call this function "at the same time", while making the critical section as short as possible?

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#### Locks

Consider the following code fragment:

```
1 int a = 0;
2
3 void f(int x) {
4    x++;
    lock();
5    a = a + x;
    unlock();
6 }
```

Where would you add lock() and unlock() calls so that multiple threads can safely call this function "at the same time", while making the critical section as short as possible?

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## Lock: Spinning vs. Blocking

#### Spinning:

- burn CPU cycles checking the lock continuously, which is wasteful
- but as soon as the lock is released by whoever had it you grab it, which is good

#### Blocking lock:

- Go to "sleep" asking for somebody to wake you up when the key's ready, which avoids being a useless CPU hog
- But this requires much more work as the OS is now involved to put you to sleep and wake you up
  - Moving your PCB from various queues, etc.

#### Rules of thumb:

- Spinning for a long time is wasteful (wasted CPU)
- Blocking for a short time is wasteful (high overhead)

#### **Deadlocks**

- Make sure you understand resource allocation graphs
  - If the "boxes" have more than "one dot": if there is a cycle, there may be a deadlock
  - If the "gray boxes" have only "one dot": if there is a cycle, there is a deadlock
- Should we do an example?
- Remember the in-class exercise with 9 locks and 2 threads?



#### **Questions?**

- Any past homework assignments or in-class examples that we should look at?
  - Perhaps examples in the IPC lecture notes for processes (output redirection, dup(), etc.)
  - The scheduling assignment?
- I have yet-another-fork example after this slide that we can look at...



## **Another fork() Example**

```
int a = 10;
int p1, p2;
p1 = fork();
if (p1 != 0) {
    a++;
    sleep(200);
    p2 = fork();
    if (p2 == 0) {
      sleep(300);
      a++;
    printf("%d\n",a);
} else {
    printf("%d\n",a);
```

What does this code print?



## **Another fork() Example**

```
int a = 10;
int p1, p2;
p1 = fork();
if (p1 != 0) {
    a++;
    sleep(200);
    p2 = fork();
    if (p2 == 0) {
     sleep(300);
      a++;
    printf("%d\n",a);
} else {
    printf("%d\n",a);
```

■ The output is 10, 11, 12