What is Concurrency?

ICS432 Concurrent and High-Performance Programming

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Disclaimer

- There is content in the two sets of lecture notes in this module that overlaps with ICS332
- If you took ICS332 last semester, bear with us (or zone out)
- But if you took it a long(er) time ago, experience shows this is useful!

Concurrency

- **Definition**: Execution of multiple "tasks" at the "same" time
- You have mostly written non-concurrent, or sequential, programs
 - At any point, you could stop the program and say exactly which execution is being executed, what the calling sequence is, what the runtime stack looks like, etc.
 - And there is a single answer to all the above for any execution of your program
- In a concurrent program, you design the program in terms of tasks, where each task as a "life of its own"
 - Each task has a specific job to do
 - Tasks may need to "talk" to each other
 - Tasks can be in different regions of the code or in the same region of the code a the same time
 - □ Tasks can be short-lived or last the whole program's execution
- A different way of thinking/programming

A brief history of concurrency (1)

- First machines were used in "single-user mode"
 - I declare: "I am going to use the machine for 2PM till 4PM"
 - I go in the special machine room and sit there for 2 hours
 - I try the punch cards that I have prepared in advance
 - I find bugs
 - I debug
 - etc.
- Extreme lack of productivity
 - During my "thinking time", this multi-million \$ machine does nothing

A brief history of concurrency (2)

- Batch Processing!
 - Instead of reserving the machine for a lapse of time to do all my activities (including debugging), I "submit" requests to a "queue"
 - The queue serves requests in order (possibly with priorities)
 - When my program fails and stops, somebody else gets the machine immediately
 - □ Great but: CPU idle during I/O!
 - And I/O takes foreeeeever

A brief history of concurrency (3)

Multi-programming (the 60's)

- Multiple programs reside in memory at once
- Made possible due to increased memory size
- Requires interrupts and memory protection
- Time-sharing (the 70's)
 - Multi-programming but rapid alternation between programs
 - Provides the illusion of programs all running simultaneously on the machine
- This is all in ICS332, and is what we have today
 - Virtual memory, fast context switching, etc.
- Eventually this has led to concurrency in user applications!
 - My application is "logically" multiple concurrent tasks
 - I can now implement it as concurrent tasks and the OS will run them simultaneously!
 - □ This is a main topic of this course

Concurrent Programs

A program consists of multiple files/modules/classes/functions

000 Terminal - vim - 98x69 job_request = (job_request_t)calloc(),sizeof(_job_request_t)); job_request>id = id; job_request>nodes; job_request>cactual_duration = actual_duration; job_request>requested_duration = requested_duration; job_request->channel_started = port_started; job_request->channel_done = port_done; job_request->channel_queue_size = port_queue_size; task = MSG_task_create("request",9,8,(void*)job_request); if (MSG_task_put(task,scheduler_host,PORT_JOB_REQUEST) != MSG_OK) { REGUEST", MSG_host_get_ xbt_assert10 heduler_host)): submit_job(pending_job_t pending_job submit_info_t* schedulers, int num_schedulers, across_clusters.algorithm.t across_cluster.algorithm, int across_cluster.algorithm.t inforc.gorithm, int across_cluster.agrorithm.t inforc.gorithm, int infractuster.agrorithm.t inforc.gorithm, int infractuster.agrorithm.t inforc.gorithm, int infractuster.agrorithm.t inforc.gorithm, int infractuster.agrorithm.t inforc.gorithm, int pert.stortes, int port.gories, int port.gories.ize) nt nun_target_schedulers=0: scheduler_info_t *target_schedulers=NULL: // Find out the list of appropriate schedulers, (ind_target_schedulers/snh_achedulers, pending_iob, preferred, across_cluster_algorithm, across_cluster_percentage, snum_target_schedulers, starget_schedulers); (nun_target_schedulers == 0) { (nun_target_schedulers == 0) { // sond requests to ends through profession or (===: {from, tonget.schedulers; i++) { submit_job_to_tonget.schedulers; i++) { submit_job_to_tonget.schedulers throm.cluster_percentage, port_storted, port_done, port_queue_size); ree(target_schedulers); d submit_job_to_target_scheduler(scheduler_info_t scheduler, pending_job_t pending_job, intro.cluster_aigorithm_t intro.cluster_aigorithm, int intra.cluster_argentiuter_percentulater_scheduleter, int port_started, int port_dore, int port_queue_size)

000

print_queue()

d print_queue(xbt_fifo_t q) (bt_fifo_iten_t iten; ob.dsscripton_t descripton; printf(stdam,,"---> Xd",xbt_fifo_size(q));

 $\label{eq:product} \begin{array}{l} print(\{\{(k), m_{1}, \dots, k\}, m_{2}, m_{2}$

Terminal - vim - 101x71

_start_job()

i start_job(job_descriptor_t jd, scheduler_bookkeeping_st *bk)

jd->stort_time = MSO_get_clock();

/# Notifu the submitter !

n_tork_t tork; int id = |0/21g; tork = 100_tork_id=200; if (100_tork_puttork,jd=2mulaiter,jd=2channel_started) != 100_t0() { ybt_assert0(a, Error while sending a jdb start notification');

ne

scheduler_init()

/* Alg-specific initialization */ witch (alg) { case FCFS: Beheduler_init_fcfs(bk):

000 Terminal - vim - 98x69 []runed_schedulers, num_pruned_schedulers, num_target_schedulers, target_schedulers); (*num_target_schedulers > 1) pending_job->flooded_across = 1; ee(pruned_schedulers): target_schedulers(int x, scheduler_info_t preferred, scheduler_info_t *schedulers, int r eduler_info_t *schedulers, int num_schedulers, *num_target_schedulers, scheduler_info_t **target_schedulers) nt i,j; nt r=0; nt done: num_target_schedulers = 0; or (i=0;i<HIN(x,num_schedulers);i++) { (i == 0) { /* first, pick the prefer for (j=0;j foun_schedulers;j++) { if (schedulers[j] == preferred) { se (r = random_integer_biased(0.num_schedulers-1); r = random_integer(0,num_schedulers-1); done = 1: done = 1; for (j=0;j(i;j)++) if ((*torget_schedulers)[j] == schedulers[n]) done = 0; if (done) break } (+nun_torget_schedulers)++; *torget_schedulers = REPLLOC(*torget_schedulers, (*mu.torget_schedulers) * sizeof(scheduler_info_t)); (*torget_schedulers)[(+nun_torget_schedulers -i)] = schedulers[n]; d find_Sx_target_schedulers(int x, scheduler_info_t preferred, scheduler_info_t *schedulers, int num_schedulers, int *num_target_schedulers, scheduler_info_t **target_schedulers) heduler_info_t *tosort: 000 Terminal - vim - 96x42 BT_LOG_NEW_DEFAULT_CATEGORY(RECEIVER, "Logging for the receiver process"); receiver(int args, char *argv[]) nt port=-1:

xbt_fifo_t fifo; dynar_t dynar;

/* Process the first argument */ if (seconf(argv[1], 'Xd', \$port) != 1) { xbt_assert1(0, 'Invalid port 'Xs' for a receiver process", argv[1]);

i get the dymon */ (({dyman = (dyman_t)xbt_dict_get_or_null(receiver_dyman_dict, (const char *)argv[2]))) { xbt_assert(d,q,"cannot (ind dyman '%s' for a receiver process',argv[2]);

/* main loop */ while(1) { int r; <u>n_task_t task</u> = NULL;

if (MSG_task_get(&task,port) != MSG_OK) {
 xbt_assert0(0,"Error while receiving a task in a receiver process");

/* Put the task at a random location in the dynam */ $r=nandom_integer(0,HRX(0,dynam_length(dynam)-1)); r=0; dynam_insert_at(dynam,dynam_length(dynam),(void*)task);$

return 1;







Concurrent Programs

- Thinking about what a concurrent program does is more difficult than for a sequential program
 - One may have to keep a mental picture of what each task is doing at all time (we try not to)
 - Questions like "While task #1 is in function f where is task #2?" are often difficult (and we try not to have their answers matter so that we don't have to ask them)
 - Two executions of the same program may not be identical
 - We'll explain this in more details
- As a result, concurrent programs are
 - Almost always more difficult to design for correctness
 - Almost always more difficult to read
 - Always more difficult to debug
- So, why do we bother at all?

Concurrency for Interactivity

- One of the oldest uses for concurrency is to make programs more interactive
- While a program is running and doing stuff, the user should still be able to interact with it
- Example:
 - What if in your Web browser you couldn't click "back" before the browser has finished loading the page you immediately realized is the wrong one?
 - What if in iTunes you couldn't look at your play list while you're playing a song because the program is busy playing the song?
- One wants to avoid the "frozen because I am working" problem as much as possible
- Let's look at a made-up example...

Designing a Concurrent GUI

- A common application of concurrent programming is for designing Graphical User Interfaces (GUIs)
- Example application
 - Say you want to write a program that renders 3-D objects on the screen
 - You have a clickable button to launch the rendering
 - But rendering takes a long time
 - You don't want the GUI to appear "frozen" while rendering the objects
 - For instance, you want the "Quit" and "Cancel" buttons to still work

Concurrent GUI?

One way to avoid the "frozen" problem without using concurrency is to write your code with breaking down a task into sub-tasks

Typically, I'll write code fragments in C/C++-like pseudo-code or in Java, without declarations, etc.

```
void render(...) {
  for (step=0; step<100; step++) {
    this.doSomeRendering(...);
    if (gui.cancelButton.clicked())
        break;
  }</pre>
```

That was (often) a bad idea

It's cumbersome:

- What if you want to do 7 tasks?
- Sprinkle "interaction checks" throughout your code will make is unreadable and annoying to maintain
- It's no always doable:
 - What if rendering is not breakable into multiple calls??
 - Perhaps you call some library that you didn't write and cannot modify
- What if some tasks have some real-time requirements?
 - e.g., you want to have an animated symbol that changes every t milliseconds but a call to doSimpleRendering() takes longer than that?

Task-based design

- Instead of thinking of your application as one task that has to juggle many things at once, you think of your application as a bunch of tasks that run concurrently
- Each task does one thing and perhaps doesn't even know that there are other tasks
- Assuming that we have a programming language that allows us to define tasks we can rewrite our application...

Concurrent GUI

In *horrible* pseudo-code:

```
renderer = new Task(render)
mousewatcher = new Task(watchmouse)
```

```
renderer.start();
mousewatcher.start();
renderer.wait_until_finished()
```

```
Renderer::render() {
   // do rendering on screen
}
```

```
Mousewatcher::watchmouse() {
   // whenever mouse clicked, kill renderer
}
```

Concurrent Tasks Abstraction

- Fortunately, our OSs support the concurrent tasks abstraction
 - After all, on our machines many programs can run simultaneously
 - So why not tasks within our programs?
- This can be done by
 - □ A special library
 - A virtual machine like the JVM
 - The Operating System
 - □ A combination of the above
- Almost all modern programming languages allow you to create "tasks" in your programs

"Simultaneous" tasks?

- Can we really have simultaneous concurrent tasks?
- There are two kinds of concurrency:

 True concurrency: two or more "things"
 happen at the same instant in time
 - False concurrency: only one thing happens at a time, but the illusion of concurrency is achieved because the OS performs rapid context switching

True/False Concurrency

- Consider a program that defines two concurrent tasks, T1 and T2
- On a single core, only one task can use the CPU
 - The concurrent tasks use false concurrency
- On a multi-core system, each task can be on a different core

The concurrent tasks use true concurrency

False Concurrency on One Core



- False concurrency between the red task, the green task, and the blue task
- The OS context-switches back and forth between the three tasks
- Because this is very fast, we have the "illusion" of simultaneous execution





green task, the grey and the blue, etc.

True/False Concurrency

- The programmer shouldn't have to care/know whether concurrency will be true or false
 - Besides the fact that true concurrency offers better performance than false concurrency
- Typically, the programmer doesn't know on which computer the program will run!
 - You have no idea how many cores your "customer" will have on their machines
- A concurrent program with 10 tasks will work on a single-core processor, a quad-core processor, a 32-core processor, etc
 - Your job as a developer is to create tasks
 - e.g., the program could easily discover that the machine it's running on has 8 cores, and thus decides to create 8 tasks
 - The job of the OS is to dispatch these tasks to the cores
 - e.g., the OS is smart enough to put each of the 8 tasks on its own core without you having to make those decisions

Performance!!!

- But wait, with true concurrency we can also go faster!!!
 - If you have to bake 2 cakes and you have 1 oven it will take you 2 hours
 - But if you have 2 ovens that can be on at the same time, it will take you only 1 hour
- This brings us to the second major reason why people want to use concurrency: compute stuff faster
- To summarize we have two motivations:
 concurrency for interactivity
 concurrency for performance

Multi-core Processors

- There have always been different hardware resources to use concurrently to increase performance
 - e.g., the disk, the network, and the CPU can all be used at the same time because they are different pieces of hardware
- But the last decades have seen the advent of multicore processors, which are now ubiquitous
- Many programs have been made concurrent so as to utilize multiple cores concurrently
 - It's become impossible to say "I am an employable software developer but I don't deal with concurrency"
- How come we have multi-core processors in the first place?

Moore's Law

In 1965, Gordon Moore (co-founder of Intel) predicted that transistor density of semiconductor chips would double roughly every 24 months (often "misquoted" as 18 months)



- □ He was right
- But, the law was often wrongly interpreted as:
 "Computers get twice as fast every 2 years"
- This wrong interpretation was true for a while, but no longer...

50-year Trend



Plot inspired from the work of Pedro Bruel, generated with data from Wikipedia [Wik21a; Wik21b].

50-year Trend



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50-year Trend



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Multi-core Chips

Constructors cannot increase clock rate further

- Power/heat issues
- They bring you multi-core processors
 - □ Multiple "low" clock rate processors on a chip
- It's really a solution to a problem, not a cool new advance
 - Even though there are many cool/interesting things about multicore processors
 - Even though writing concurrent code is cool/interesting, as we'll see in this course
- But most users/programmers would rather have a 100GHz core than 50 2GHz cores
 - □ In which case we would not need to write concurrent programs
 - When given the choice, if you can get by without concurrency, you're likely better off (until you can't avoid concurrency anymore)
 - i.e., in general no compiler will nicely take your sequential app and magically transform it into an efficient multi-threaded app

So, Multi-Core = High Performance?

- A big question is: how much performance benefit can we really get from concurrency?
- It's a difficult question because the answer
 - depends on the application
 - depends on the computer
 - depends on the language / operating system
- In some cases, it's very easy to achieve great performance via concurrency
- In others, it's very difficult
- We'll be exposed to this in this course

Take-away

- Concurrency is about structuring your programs as sets of tasks
- Typically done for interactivity and/or for performance
- Concurrency is supported by programming languages, by OSes, and by the hardware
- Issues for programmers:
 - Correctness (we'll see this can be a tough one)
 - Performance (sometimes easy, sometimes not)

Task-based Thinking

- From now on, you should begin writing code with concurrency in mind (even if the code is not concurrent right away)
- You currently think of your programs as sets of classes/objects
 Or data structures and functions
- But now, you also have to think of your programs as sets of tasks
 And these tasks operate on the objects
- This requires a little bit of mental adjustment, and our programming assignments will help making that adjustment
 - □ You have been exposed to this a little bit in ICS332, which will help too
- Thinking of programs as "sets of tasks" will become second nature soon enough
 - And is very natural for many applications actually
 - e.g., each tab in your Web browser is a task
- When I write a sequential program, I typically think of it as a concurrent program that happens to have a single task

Conclusion

- In the next set of lecture notes, we'll look at processes and threads
 - □ What they are (super quick ICS332 "review")
 - How we can use them to make an application concurrent (beyond ICS332)